

Choice from Finite Sets: A Topos View

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Introduction

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Mostowski-Gauntt Permutation Group Criteria

Theorem (Local)

For any natural number n and any finite set Z of natural numbers, the following are equivalent.

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Fix a group *G*.

G-Set: topos of sets with (left) action of *G* and functions that commute with the actions.

More Examples

M -sets for any monoid M

Kripke structures for intuitionistic logic

Presheaves on any small category

Sheaves on any topological space (or any site)

Realizability topoi

Facts about \mathbf{Set}^I

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Notation: Abbreviate “In topos \mathcal{E} , every n -element set has a point” as $EP(n, \mathcal{E})$. And $EP(Z, \mathcal{E})$ means $EP(z, \mathcal{E})$ for all z in the finite set Z .

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So $EP(n, G\text{-Set})$ says that every action of G on n has a fixed point. And $EP(Z, G\text{-Set})$ says the same for actions of G on any $z \in Z$.

Local Mostowski-Gauntt in Topos Language

Theorem

For any natural number n and any finite set Z of natural numbers, the following are equivalent.

- *For every topos \mathcal{E} of the form **Model of ZF/I**, $EP(Z, \mathcal{E})$ implies $EP(n, \mathcal{E})$.*
- *For every topos \mathcal{E} of the form **G-Set**, $EP(Z, \mathcal{E})$ implies $EP(n, \mathcal{E})$.*

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What other sorts of topoi does this equivalence apply to?

General Local Theorem

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General Global Theorem

Theorem

For any natural number n and any finite set Z of natural numbers, the following are equivalent.

- *ZF (or ZFA) proves that, if $C(Z, I)$ for all I then $C(n, I)$ for all I .*
- *For every topos \mathcal{E} of the form **Model of ZF**/ I , if $EP(Z, \mathcal{E}/F)$ for all finite decidable objects F of \mathcal{E} , then $EP(n, \mathcal{E})$.*
- *Every group that acts without fixed points on n has a subgroup that acts without fixed points on some $z \in Z$.*
- *For every topos \mathcal{E} of the form **G-Set**, if $EP(Z, \mathcal{E}/F)$ for all finite decidable objects F of \mathcal{E} , then $EP(Z, \mathcal{E})$.*
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